

Reverse Mathematics and Marriage Problems

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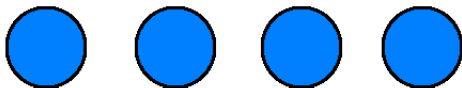
Agenda

- ▶ I: Marriage Problems
- ▶ I: Previous Results
- ▶ I: New Results
- ▶ II: Reverse Mathematics

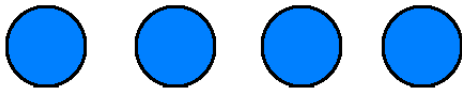
I: Marriage Problems

Marriage Problems

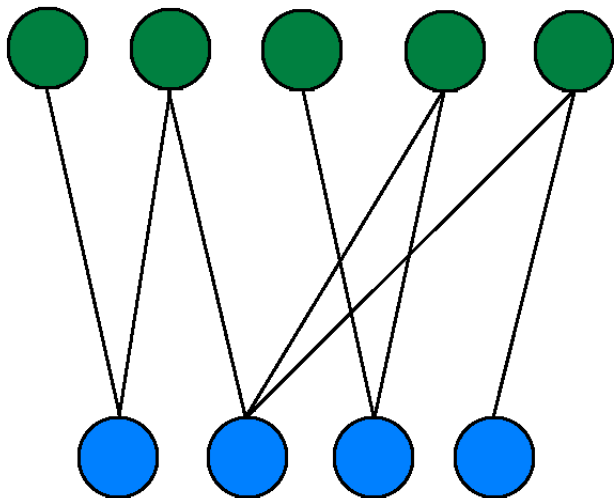
Marriage Problems



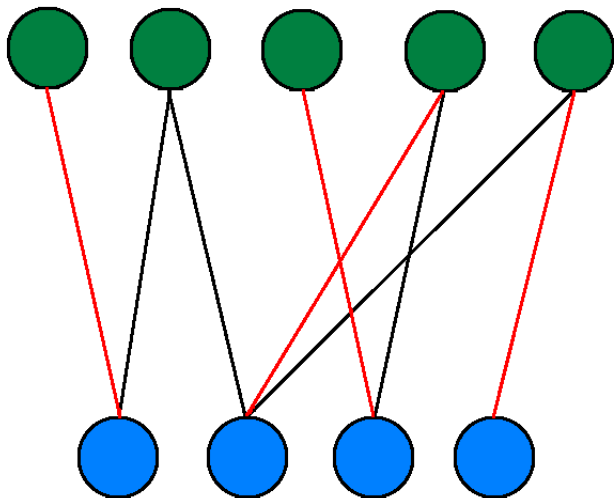
Marriage Problems



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Some Notation

A marriage problem M consists of three sets B , G and R .

B is the set of boys,

G is the set of girls, and

R is the relation between the boys and girls.

$R \subset B \times G$ where $(b, g) \in R$ means “ b knows g ”.

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$G_M(b)$ denotes the set of girls b knows relative to the relation in M .

Some More Notation

A *solution* to $M = (B, G, R)$ is an injection

$$f : B \rightarrow G$$

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M is a:

finite marriage problem if $|B|$ is finite.

infinite marriage problem if $|B|$ is not finite.

bounded marriage problem if there is a function $h : B \rightarrow G$ so that for each $b \in B$, $G(b) \subseteq \{0, 1, \dots, h(b)\}$.

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Theorem

If $M = (B, G, R)$ is a finite marriage problem such that $|G(B_0)| \geq |B_0|$ for every $B_0 \subset B$, then M has a solution.

Due to Philip Hall.

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In the finite case, we found the following necessary and sufficient condition.

Theorem

(RCA₀) *If $M = (B, G, R)$ is a finite marriage problem with n boys and a unique solution f , then there is an enumeration of the boys $\langle b_i \rangle_{i \leq n}$ such that for every $1 \leq m \leq n$, $|G(\{b_1, b_2, \dots, b_m\})| = m$.*

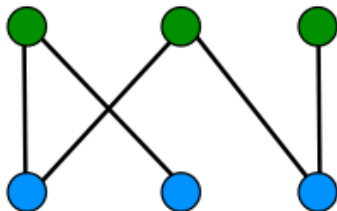
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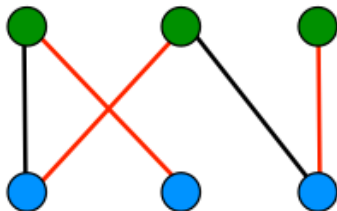
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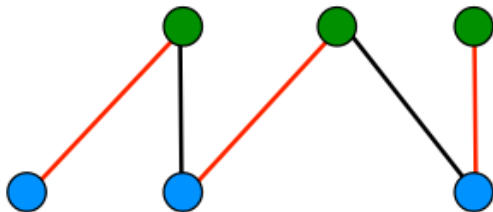
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Sketch of the proof

Lemma

(RCA_0) *If $M = (B, G, R)$ is a finite marriage problem with a unique solution f , then some boy knows exactly one girl.*

Sketch of the proof

Proof: Suppose we have $M = (B, G, R)$ as stated above with some initial enumeration of B . Apply the lemma and let b_1 be the first boy such that $|G(b_1)| = 1$.

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Define $M_2 = (B - \{b_1\}, G - G(b_1), R_2)$. Because M has a unique solution, M_2 has a unique solution, namely the restriction of f to the sets of M_2 . Apply the lemma once more and let b_2 be the first boy in $B - \{b_1\}$ such that $|G_{M_2}(b_2)| = 1$.

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Continuing this process inductively yields the j^{th} boy in our desired enumeration from

$$M_j = (B - \{b_1, b_2, \dots, b_{j-1}\}, G - G(b_1, b_2, \dots, b_{j-1}), R_j).$$

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After the n^{th} iteration we have (b_1, b_2, \dots, b_n) where for every $1 \leq m \leq n$, $|G(\{b_1, b_2, \dots, b_m\})| = m$. ■

Generalizing this result

The statement regarding finite marriage problems with unique solutions can be generalized to the *infinite* case. Paralleling the previous work we have:

Theorem

If $M = (B, G, R)$ is an infinite marriage problem with a unique solution f , then there is an enumeration of the boys $\langle b_i \rangle_{i \geq 1}$ such that for every $n \geq 1$, $|G(\{b_1, b_2, \dots, b_n\})| = n$.

II: Reverse Mathematics

Reverse Mathematics

Reverse mathematics is the subfield of mathematical logic dedicated to classifying the **logical strength** of mathematical theorems.

This is done by proving theorems equivalent to a hierarchy of axioms over a weak base axiom system.

$$\text{RCA}_0 \quad \text{WKL}_0 \quad \text{ACA}_0 \quad \text{ATR}_0 \quad \Pi_1^1 - \text{CA}_0$$

RCA_0 proves the *intermediate value theorem* and the *uncountability of \mathbb{R}* .

RCA_0 does **not** prove the *existence of Riemann integrals*.

Equivalences

Theorem

The following are provable in RCA_0 .

- (i) $\text{WKL}_0 \iff$ *For every continuous function $f(x)$ on a closed and bounded interval $a \leq x \leq b$, the Riemann integral $\int_a^b f(x)dx$ exists and is finite. (Simpson)*
- (ii) $\text{ACA}_0 \iff$ *For all one-to-one functions $f : \mathbb{N} \rightarrow \mathbb{N}$ there exists a set $X \subseteq \mathbb{N}$ such that $\text{Ran}(f) = X$. (Simpson)*
- (iii) $\text{ATR}_0 \iff$ *Any two well orderings are comparable. (Friedman)*
- (iv) $\Pi_1^1 - \text{CA}_0 \iff$ *The Cantor/Bendixson theorem for $\mathbb{N}^{\mathbb{N}}$: Every closed set in $\mathbb{N}^{\mathbb{N}}$ is the union of a perfect closed set and a countable set. (Simpson)*

Marriage Theorems and Reverse Mathematics

Jeff Hirst proved the following equivalence results:

Theorem

(RCA_0) *If $M = (B, G, R)$ is a finite marriage problem such that $|G(B_0)| \geq |B_0|$ for every $B_0 \subset B$, then M has a solution.*

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(RCA_0) The following are equivalent:

- 1 ACA_0
- 2 If $M = (B, G, R)$ is an infinite marriage problem such that $|G(B_0)| \geq |B_0|$ for every $B_0 \subset B$, then M has a solution.

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- 2 If $M = (B, G, R)$ is a bounded marriage problem such that $|G(B_0)| \geq |B_0|$ for every $B_0 \subset B$, then M has a solution.

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Our new results echoed the previous work:

Theorem

(RCA₀) *If $M = (B, G, R)$ is a finite marriage problem with n boys a unique solution f , then there is an enumeration of the boys $\langle b_i \rangle_{i \leq n}$ such that for every $1 \leq m \leq n$, $|G(\{b_1, b_2, \dots, b_m\})| = m$.*

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Future Work

- ▶ Marriage problems with any fixed finite number of solutions.
- ▶ “Entangled societies”

References

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